# Contingency tables and hypergeometric polynomials associated to hyperplane arrangements

Nobuki Takayama, joint work with Y.Tachibana, Y.Goto, T.Koyama

## 2018/06/11

- Yoshihito Tachibana, Yoshiaki Goto, Tamio Koyama, Nobuki Takayama, Holonomic Gradient Method for Two Way Contingency Tables, arxiv:1803.04170
- Y.Goto, K.Matsumoto, Pfaffian equations and contiguity relations of the hypergeometric function of type (k+1,k+n+2) and their applications, arxiv:1602.01637

$$I = (I_1, \dots, I_{k+1}) \in \mathbf{Z}_{\geq 0}^{k+1}, \ J = (J_1, \dots, J_{n+1}) \in \mathbf{Z}_{\geq 0}^{n+1},$$
  

$$\sum I_i = \sum J_j.$$
  

$$p = (p_{ij})$$
  

$$Z(I, J; p) = C.T. \prod_{j=1}^{n+1} \left(\sum_{i=1}^{k+1} p_{ij} t_i\right)^{J_j} t^{-I}$$
(1)

 $t_1 = 1$ ,  $t^{-l} = \prod_{i=1}^{k+1} t_i^{-l_i}$ . Note\* that

$$Z(I, J; p) = J! \sum \frac{p^u}{u!}$$

where  $\sum_{i} u_{ij} = J_j$  (column sum is J),  $\sum_{j} u_{ij} = I_i$  (raw sum is I)<sup>†</sup>. Z is the normalizing constant (partition function) of a distribution.

<sup>\*</sup>C.T. is the constant term w.r.t.  $t. J! = \prod_i J_j!$ 

<sup>&</sup>lt;sup>†</sup>We denote these conditions or the left hand sides of them by ©or 🕮 🖉 🕤

Goal 1: Evaluate numerically Z and its derivatives efficiently and accurately  $\ddagger$ .

Motivation from statistics: 2 way contingency table:

 $(k+1) \times (n+1)$  matrix with  $Z_{\geq 0}$  entries.

	acetaminophen	diclofenac sodium	mefenamic acid
death	4	7	2
survival	32	5	6

$$P(U_{ij} = u_{ij}) = rac{\exp(-
ho_{ij})
ho_{ij}^{u_{ij}}}{u_{ij}!}$$

The conditional probability  ${}^{\$}$  when the row and column sums are fixed to  $I,\ J$  is

$$P\left(U=u \mid \sum_{j} U_{ij} = I_i, \sum_{i} U_{ij} = J_j\right) = \frac{p^u/u!}{Z(I, J; p)}$$

<sup>‡</sup>When I = (4, 14, 5, 2, 1), J = (10, 6, 5, 2, 3), there are 229, 174 terms. <sup>§</sup> $U_{ij}$  is a random variable of the Poisson distribution A = A = A = A

#### References on contingency tables (MSC2010: 62H17).

Takayuki Hibi *Editor* 



# **Gröbner Bases**

Statistics and Software Systems



$$E[U_{ij}|\odot] = \sum_{\odot} \frac{u_{ij}p^u/u!}{Z(I,J;p)} = p_{ij}\frac{\partial}{\partial p_{ij}}\log Z$$
(2)

#### Proposition

 $E[U_{ij}|\odot_U]$  is invariant by the torus action  $p_{ij} \mapsto p_{ij}p_ip'_j$ ,  $p_i, p'_j \in \mathbf{R}_{>0}$ .

Theorem

$$\mathbf{R}_{>0}^{(k+1)(n+1)}/\sim \ni (p_{ij})\mapsto E[U_{ij}|\odot] \in \operatorname{relint}\operatorname{New}(Z)$$
(3) is an isomorphism ¶.

Goal 2: Find the inverse image numerically  $\parallel$ .

 $<sup>\</sup>P_{\sim}$  is the equivalence relation w.r.t. the torus action. This theorem is a special case of Th. 1 of N.Takayama, S.Kuriki, A.Takemura, A-Hypergeometric Distributions and Newton Polytopes, Advances in Applied Mathematics 99 (2018) 109–133.

Let us explain the idea of our method \*\* for  $2 \times 2$  case.

$$\bar{u} = \begin{pmatrix} J_1 & 0 \\ J_2 - I_1 & I_2 \end{pmatrix}.$$

$$Z = \frac{p^{\bar{u}}}{\bar{u}!} {}_2F_1 \left( -J_1, -I_2, J_2 - I_2 + 1; \frac{p_{12}p_{21}}{p_{11}p_{22}} \right)$$
(4)
$$f(a) = {}_2F_1(a, b, c; x) = \sum_{k=0}^{\infty} \frac{(a)_k(b)_k}{(c)_k(1)_k} x^k,$$
(a)<sub>k</sub> = a(a + 1) ··· (a + k - 1). F(a) = (f(a), xdf/dx(a))^T.
$$F(a) = (E + A(a)/a)^{-1}F(a + 1),$$
(5)
where  $A(a) = \begin{pmatrix} 0 & 1 \\ abx/(1-x) & (ax + bx - c + 1)/(1-x) \end{pmatrix}.$ 

$$F(a) = M(a)M(a+1) \cdots M(-2)F(-1), \quad M(a) = (E + A(a)/a)^{-1}.$$

"factorial" of contiguity relation (5).

\*\*holonomic gradient method, HGM

 $L_{j} = \tilde{p}_{j} \cdot t \text{ where } \tilde{p}_{j} \text{ is the } j\text{-th column vector }^{\dagger\dagger} \text{ of } \tilde{p}. \ \alpha_{j} \in \mathbf{C} \setminus \mathbf{Z},$   $\sum_{j=1}^{k+n+2} \alpha_{j} = 0.$   $\nabla = d_{t} + \sum_{j} \alpha_{j} d_{t} \log L_{j} \qquad (6)$   $\tilde{P} = \{\tilde{p} \mid \text{ any } (k+1) \times (n+1) \text{ minor of } \tilde{p} \neq 0\}$ 

$$\mathcal{T}_{p} = \{t' \in \mathbf{C}^{k} \mid L_{j}(p; t) \neq 0 \text{ for all } j.\}, \quad p \in \tilde{P}$$

 ${}^{\dagger\dagger}L_1 = t_1 = 1, L_2 = t_2, \dots, L_{k+n+2} = \sum_i p_{i,n+1} t_i \quad \langle \square \rangle \quad (\square ) \quad \langle \square \rangle \quad \langle \square \cap \cap \cap \quad \langle \square \cap \cap \cap \quad \langle \square \cap \cap \cap \cap \quad \langle \square \cap \cap \cap \cap \quad (\square ) \quad ($ 

$$\mathcal{J} = \{(j_1, \dots, j_{k+1}) \mid 1 \leq j_1 < j_2 < \dots < j_{k+1} \leq k+n+2\}$$

$$q\mathcal{J}_p = \{J \in \mathcal{J} \mid q \notin J, p \in J\}$$

$$\varphi\langle J \rangle = d_t \log(L_{j_2}/L_{j_1}) \wedge d_t \log(L_{j_3}/L_{j_1}) \wedge \dots \wedge d_t \log(L_{j_{k+1}}/L_{j_1}) \quad (7)$$
Theorem (Goto, Matsumoto(2016) \*, contiguity relation)
$$Put \ F = (\varphi\langle J \rangle \mid \quad J \in {}_{k+n+2}\mathcal{J}_1) \text{ and assume } \tilde{p} \in \tilde{P} \text{ Then,}$$

$$L_i F \equiv (CP_i^{-1}D_iQ_iC^{-1}) F \quad \text{in} \quad H^k(\Omega^{\bullet}(T_p), \nabla) \quad (8)$$

where  $C, P_i, Q_i$  are intersection matrices <sup>†</sup> among  $\varphi \langle J \rangle$  and  $D_i$  is a diagonal matrix with rational function entries of p.

\*Y.Goto, K.Matsumoto, Pfaffian equations and contiguity relations of the hypergeometric function of type (k+1,k+n+2) and their applications, arxiv:1602.01637, to appear in Funkcialaj Ekvacioj.

 ${}^{\dagger}H^{k}(\mathcal{E}_{0}^{\bullet}(T_{p}), \nabla^{\vee}) \times H^{k}(\Omega^{\bullet}(T_{p}), \nabla) \to \mathbf{C} \text{ is called the intersection form.}$ Note that  $H^{k}(\mathcal{E}_{0}^{\bullet}(T_{p}), \nabla^{\vee}) \simeq H^{k}(\Omega^{\bullet}(T_{p}), \nabla^{\vee})$ 

Integrating the both sides of (8) with  $\int_{\Delta} \prod L_j^{\alpha_j}$ , we have a contiguity relation for Z.

#### Theorem

Fix k. Then the complexity of constructing the contiguity relation is  $O(n^{3(k+1)})$ .

How do we evaluate efficiently  $M(a)M(a-1)\cdots M(-1)$ ?  $\Rightarrow$  the modular method in computer algebra; evaluate in  $\mathbb{Z}/s\mathbb{Z}$  for

several prime numbers s and reconstruct the answer in  $\mathbf{Q}$  by the Chinese remainder theorem.

#### Theorem

Let n be the number of linear transformations and put  $r = {\binom{k+n}{k}}^{\ddagger}$ . The complexity of the modular method is  $\max(O(|J|), O(r)), |J| = \sum J_j.$ (Numerical evidences.)

<sup>&</sup>lt;sup>‡</sup>the rank of the twisted cohomology group.  $(\Box \rightarrow \langle \Box \rangle \land \Xi \rightarrow \langle \Xi \rangle \land \Xi \rightarrow \exists \neg \neg \neg \neg \neg$ 

Have we solved two goals?  $\Rightarrow$  Not completely. We have assumed that  $\tilde{p}\in\tilde{P}$  §.

### Proposition

Let  $\beta_1$  be the total degree of Z and L a generic line in p-space. If we evaluate  $E[U_{ij}]^{\P}$  at  $2\beta_1$  points  $p \in \mathbf{R}_{>0}^{(k+1)\times(n+1)}$  on a line L, then the exact value of  $E[U_{ij}]$  can be obtained at any point on L. However, this method is not efficient  $\Rightarrow$  open questions  $\parallel$  for hyperplane arrangements of the case that some of  $p_{ij} = 0$ .



Figure:  $V(t_2t_3(p_{21}t_2+p_{31}t_3)\prod_{j=2}^3(p_{1j}+p_{2j}t_2+p_{3j}t_3))$ 

<sup>§</sup>This is the condition that hyperplane arranement is in a generic position. <sup>¶</sup>We denote  $E[U_{ij}|\odot_U]$  by  $E[U_{ij}]$ . <sup>∥</sup>Y.Goto, 1805.01714 This book will help to solve the open question of constructing contiguity relations efficiently for any hyperplane arrangement.



What is the space

$${\sf R}^{(k+1)(n+1)}_{\geq 0}/\sim$$

It is not a manifold!

1. Algraic geometry: Related to the Chow quotient by M.Kapranov (1992).

2. Measure theoretic (statistic).  $U_{ij}: \Omega_{ij} \rightarrow \mathbf{Z}, P(U_{ij} = u_{ij}) = \exp(-\theta_{ij})\theta_{ij}^{u_{ij}}/u_{ij}!.$  $\Omega = \prod \Omega_{ij} \times \Theta, \Theta = \{(\theta_{ij}) | \theta_{ij} \in \mathbf{R}_{\geq 0}\}.$ 

$$\mathcal{O} = \sigma \left( \odot_{U}, \frac{\theta_{ij}\theta_{k\ell}}{\theta_{i\ell}\theta_{kj}}, Z_{ij}(\theta) \right).^{**}$$
(9)

where  $Z_{ij}(c) = 1$  when  $c_{ij} > 0$  and = 0 when  $c_{ij} = 0$ .

Theorem

$$E[X|\sigma(\textcircled{o}_U,\theta)] = E[X|\sigma(\textcircled{o}_\theta,\mathcal{O})] = E[X|\mathcal{O}]$$

for any  $X \in \mathcal{L}^1(\sigma(U))$  <sup>††</sup>.

Categorial data for all:

Bed time $\setminus$ Hours slept	less than 6 hour	6–7	more than 7 hours			
Before 24	1	6	123			
24–25	3	22	145			
After 25	86	91	176			
Categorical data for males $\begin{pmatrix} 1 & 2 & 28 \\ 0 & 4 & 47 \\ 35 & 32 & 71 \end{pmatrix}$ .						
Categorical data for females $\begin{pmatrix} 0 & 4 & 95 \\ 3 & 18 & 98 \\ 51 & 59 & 105 \end{pmatrix}$ .						
/0.458167657900967 1 <u>6.25676090279981</u> \						
CMLE for males:	0 1	<u>5.252</u>	<u>00491199345</u> .			
	1 1		1 /			
CMLE for females:		\				
	<u>13.27147737376</u>	$\overline{57}$				
0.193351042187373 1	<u>3.048725861552</u>	<u>91</u> ].				
$\begin{pmatrix} 1 & 1 \end{pmatrix}$	1	< <b>f</b> > <	● ・ < ≧ > < ≧ > < ≧ > < ≧ > < ○ へ () 13 / 15			

	acetaminophen	diclofenac sodium	mefenamic acid				
death	4	7	2				
survival	32	5	6				
±‡							
CMLE: $\begin{pmatrix} 1 & \frac{10.5557279737263}{1} & 2.62096714359908 \\ 1 & 1 & 1 \end{pmatrix}$ . Generalized odds ratios: $\begin{pmatrix} 1 & \frac{11.2}{1} & 2.66666666666666667 \\ 1 & 1 & 1 \end{pmatrix}$ . $11.2 = \frac{32 \times 7}{4 \times 5}$ .							

<sup>&</sup>lt;sup>‡‡</sup>Data of the previous page

https://cran.r-project.org/web/packages/LearnBayes/index.html. Data of this papge https://www.pmda.go.jp/files/000148557.pdf =>>

Summary:

- 1. Exact numerical evaluation of the hypergeometric polynomial Z can be done efficiently with contiguity relations and the modular method.
- 2. It has applications to conditional maximal likelihood estimation (CMLE) for two way contingency tables.

Future challenge: Each column is death month, each row is birth month \*.

1	0	0	0	1	2	0	0	1	0	1	0
1	0	0	1	0	0	0	0	0	1	0	2
1	0	0	0	2	1	0	0	0	0	0	1
3	0	2	0	0	0	1	0	1	3	1	1
2	1	1	1	1	1	1	1	1	1	1	0
2	0	0	0	1	0	0	0	0	0	0	0
2	0	2	1	0	0	0	0	1	1	1	2
0	0	0	3	0	0	1	0	0	1	0	2
0	0	0	1	1	0	0	0	0	0	1	0
1	1	0	2	0	0	1	0	0	1	1	0
0	1	1	1	2	0	0	2	0	1	1	0
0	1	1	0	0	0	1	0	0	0	0	0

\*Diaconis, Sturmfels (1998), Algebraic algorithms for sampling from conditional distributions. Andrews, Herzberg (1985), Data, Springer, page 429.